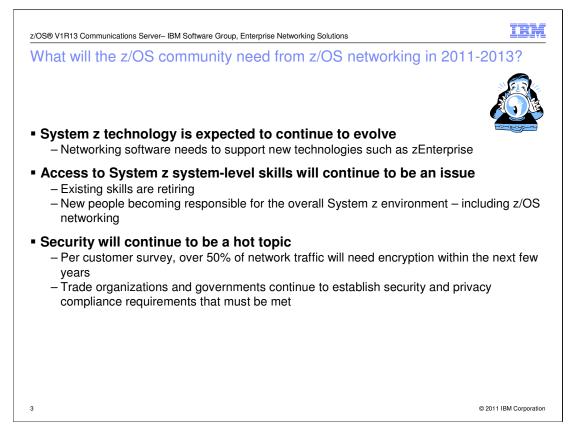


This presentation provides an overview of the new functions in z/OS V1R13 Communications Server.

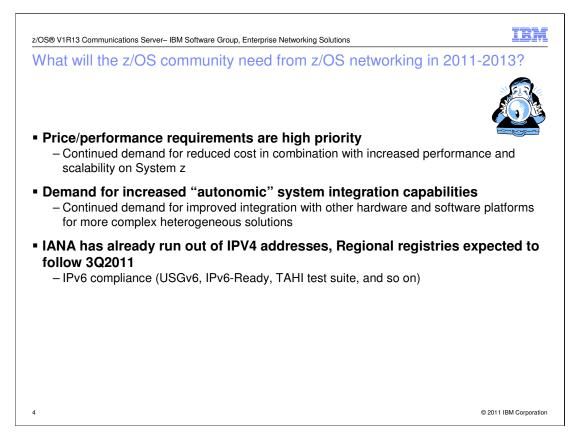
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Here is the Agenda. The first topic is a brief introduction.



In the near future, networking software needs to provide support for new System z technologies such as zEnterprise. In addition, System z skills will continue to be an issue as experienced people retire. Follow the IBM Academic Initiative

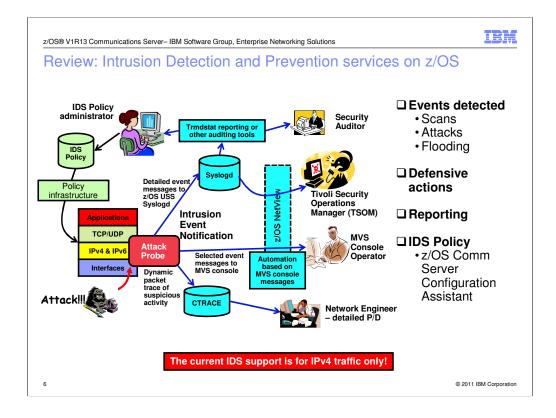
https://www.ibm.com/developerworks/university/academicinitiative/. Security has been and continues to be a hot issue facing the z/OS community.



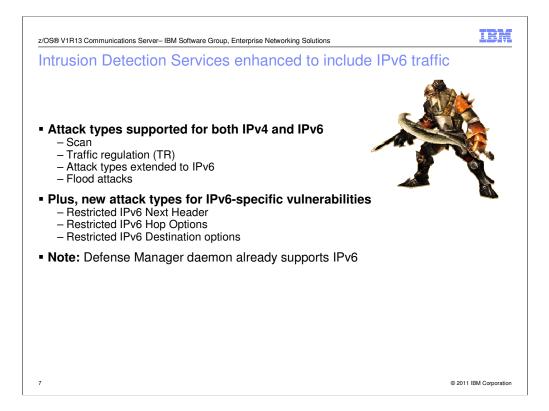
There is a continued demand for reduced costs associated with System z and for autonomic integration. IPv6 utilization is expected to increase as the Internet Assigned Numbers Authority (IANA) has already run out of IPv4 addresses. Regional registries are expected to run out of IPv4 addresses in 2011 as well.

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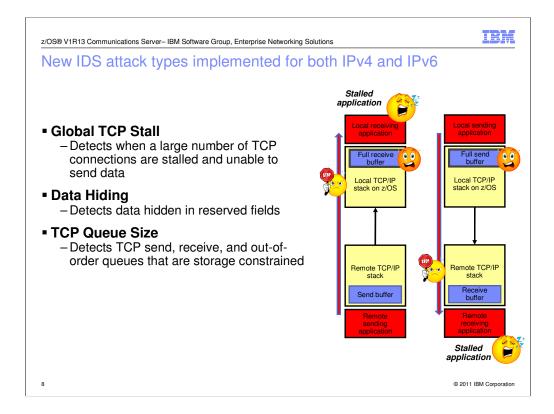
The next topic is Security enhancements.



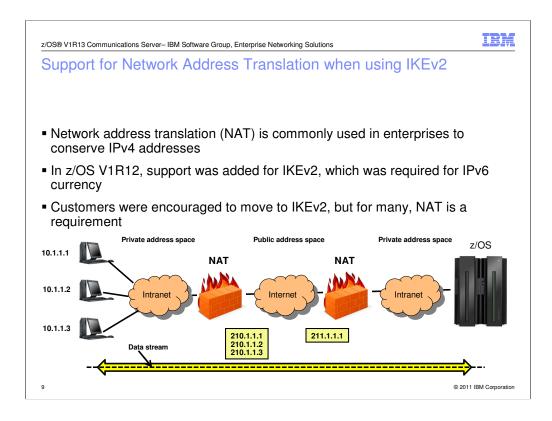
Intrusion detection services (IDS) supports detection of scans, attacks and flooding. IDS can be configured to discard packets or limit connections. Events can be logged to syslogd, to the MVS console, to IDS packet trace and to Tivoli Security Operations Manager (TSOM). IDS is configured using policy and is supported by the Configuration Assistant. Currently IDS is only supported for IPv4 traffic.



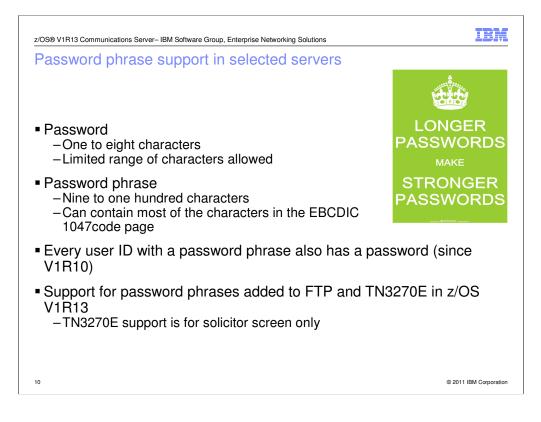
The existing event detection for scans, traffic regulation, attacks and flood attacks is enhanced to include IPv6 traffic. TCP and UDP scan event rules now support IPv6 traffic. The ICMP scan event rule is unchanged and a corresponding ICMPv6 scan event rule is added. Scan exclusion lists can now include IPv6 addresses. TCP Traffic Regulation (TR) is enhanced to monitor IPv4 and IPv6 connection requests and UDP TR is enhanced to monitor IPv4 and IPv6 packets. The malformed packet event, UDP perpetual echo, and ICMP redirect restrictions attacks are extended to IPv6. IPv6 packets can be dropped due to malformed headers, options, or values. The SYN flood and interface flood attacks are extended to IPv6 as well. New attacks added for IPv6-specific vulnerabilities include Restricted IPv6 Next Header, Restricted IPv6 Hop Options, and Restricted IPv4 Destination options. The Defense Manager daemon already supports IPv6.



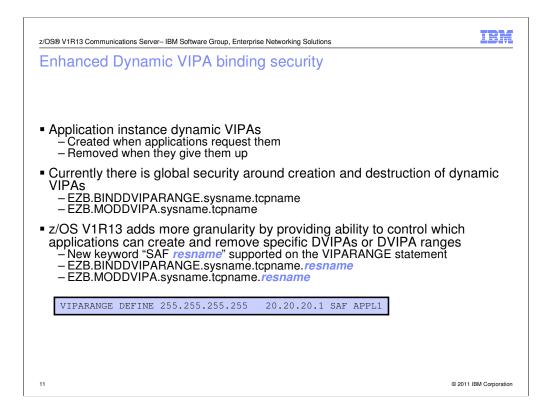
IDS implemented three new attack types for IPv4 and IPv6. The global TCP stall attack type prevents an attacker from creating connections with zero window sizes and keeping them open indefinitely. The data hiding attack type prevents an attacker from hiding data in reserved fields. This can be PadN options in IPv6 and reserved fields in IPv4 headers. The TCP queue size attack type helps you manage the amount of storage TCP can take up for the queues holding sent and received data. For example, out of order packets awaiting re-sequencing. It provides user control over storage constraint availability improvements added in z/OS V1R11. This helps avoid TCP storage constraint situations.



NAT is commonly used to conserve IPv4 addresses. IKEv2 support was added in V1R12 and supports both IPv4 and IPv6. Network address translation (NAT) is now supported when using IKEv2. You can now migrate from IKEv1 to IKEv2 if you are using NAT.



The FTP and TN3270E servers have been updated to support password phrases. Passwords are one to eight characters in length and have a limited range of characters allowed. For example, a space is not allowed in a password. Password phrases extend the length to 100 characters and support most of the characters in the 1047 code page.



Application instance dynamic VIPAs are virtual IP addresses that are created when applications request them and removed when they give them up. They provide improved availability. For example a dynamic VIPA can move around in the sysplex, following the application when it moves, so clients are uninterrupted.

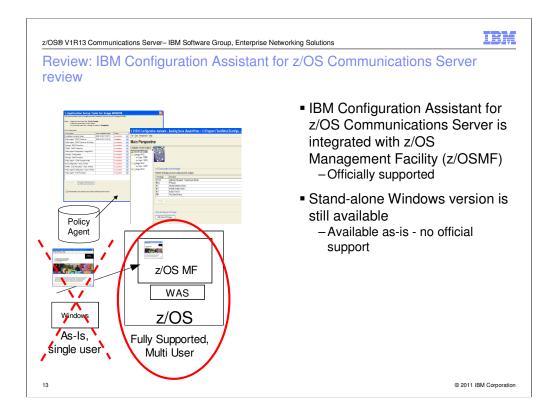
Currently there is global security around creation and destruction of dynamic VIPAs. An application can be permitted to create and destroy all dynamic VIPAs. An application permitted to

EZB.BINDDVIPARANGE.sysname.tcpname can bind to and remove all VIPARANGE defined DVIPAs. Similarly an application permitted to EZB.MODDVIPA.sysname.tcpname can issue MODDVIPA or SIOCSVIPA to create and remove all VIPARANGE defined DVIPAs.

z/OS V1R13 adds more granularity by providing the ability to control which applications can create and remove specific DVIPAs or DVIPA ranges. This allows an application to create/remove its own DVIPAs but prevent it from interfering with other applications' ranges. A new keyword "SAF resname" is supported on the VIPARANGE statement. This identifies the resource profiles to use when creating or removing DVIPAs for the VIPARANGE statement. If the SAF keyword is not present, the existing profiles are used. In the example, to bind to 20.20.20.1, the application must be permitted to EZB.BINDDVIPARANGE.sysname.tcpname.APPL1. To issue MODDVIPA 20.20.20.1, the application must be permitted to EZB.MODDVIPA.sysname.tcpname.APPL1.

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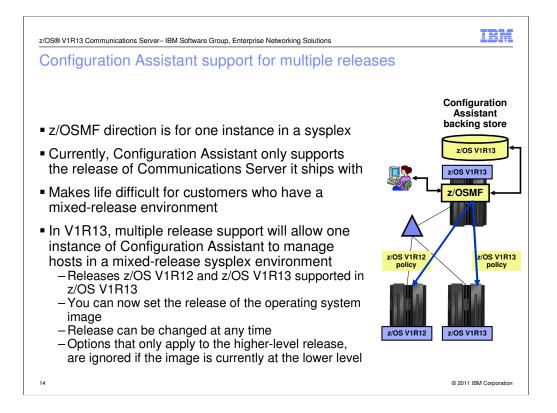
The next topic is Simplification enhancements.



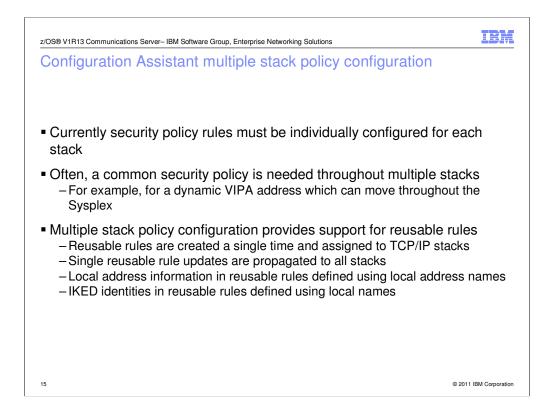
As of z/OS V1R11, IBM Configuration Assistant for z/OS Communications Server is integrated with z/OS Management Facility (z/OSMF). The z/OSMF version is integrated into the product and runs on z/OS. It is officially supported.

The stand-alone Windows version is still available, but is made available as-is, without any official support. It can be found at

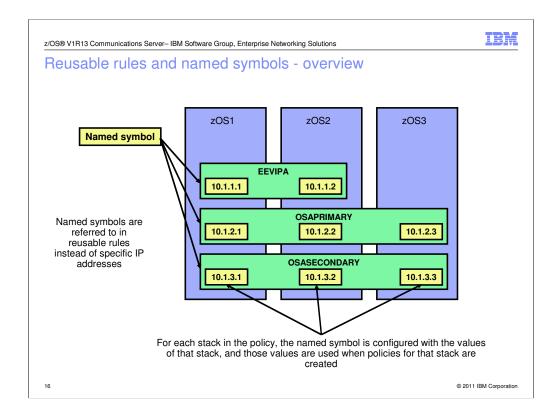
http://www.ibm.com/support/docview.wss?rs=852&context=SSSN3L&dc=D400&uid =swg24013160&loc=en\_US&cs=UTF-8&lang=en&rss=ct852other or http://tinyurl.com/cgoqsa.



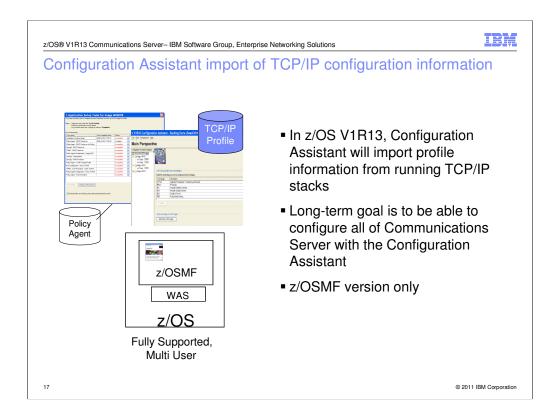
The z/OSMF direction is for a single instance of z/OSMF in a sysplex. For customers who have a mixed-release environment, this means that one instance has to manage multiple releases. Currently, the version of Configuration Assistant that ships with a release of Communications Server can only update that version. Multiple release support will allow one instance of Configuration Assistant to manage hosts in a mixed-release sysplex environment. Releases z/OS V1R12 and z/OS V1R13 are supported in z/OS V1R13. When creating a new operating system image, you can now set the release of that image. The release of the operating system image can be changed at any time. This facilitates migration of images. Options that only apply to the higher-level release are ignored if the image is currently at the lower level. For example, new IDS attack types are ignored if the z/OS image is z/OS V1R12, but included if the image is z/OS V1R13.



Currently security policy rules must be individually configured for each stack in a Sysplex or other grouping. Many times, a common security policy is needed on each stack. For example, for a dynamic VIPA address which can move throughout the Sysplex. Currently, you must configure the rule for it on every host it can reside on. With z/OS V1R13, Sysplex-wide policy configuration allows for a rule to be configured once for a dynamic VIPA and pushed to every host it can reside on. This supports the Single System Image view for configuring Sysplex. Reusable rules are created a single time and assigned to TCP/IP stacks. If a reusable rule needs to be updated, only a single rule needs to be modified and the changes are propagated to all stacks. Local address information in a reusable rule is defined by way of local address names (a user-chosen symbol). Each stack that uses the reusable rule, defines the actual value of that symbol. IKED identities in reusable rules are also defined using local names.



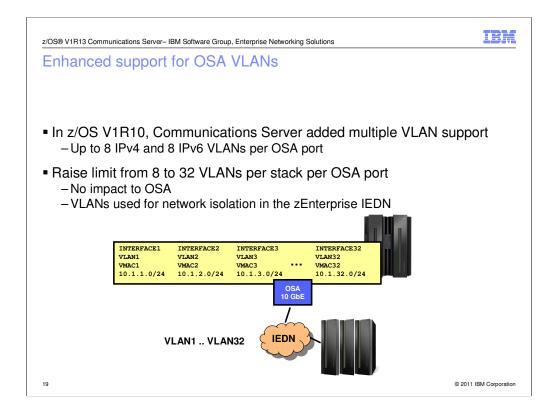
This diagram shows how reusable rules and named symbols can be used in the policy configuration for three different stacks. Reusable rules are created for all three stacks which contain named symbols. For each stack, the named symbols are configured with the values for the stack. Some of the named symbols can be dynamically configured by the Configuration Assistant.



For z/OS V1R13, Configuration Assistant will import profile information from running TCP/IP stacks. This information can be used to help develop policy configuration. For example, you can discover home IP addresses and it will suggest address groups. The long-term goal is to be able to configure all of Communications Server with the Configuration Assistant. The ability to import TCP/IP configuration information is the first step in that direction. This support is only provided for the z/OSMF version of the Configuration Assistant – support is not provided in the as-is Windows version.

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The next topic is Economics and Platform Efficiency enhancements.

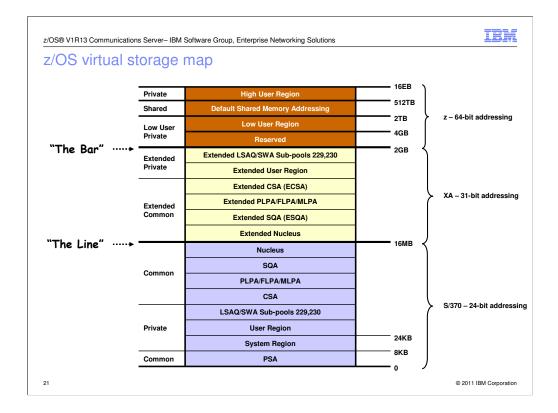


In z/OS V1R10, Communications Server added multiple VLAN support. Up to eight IPv4 and eight IPv6 VLANs are allowed per OSA port. A separate INTERFACE statement and data device are required per VLAN. The value of eight is a z/OS Communications Server software limitation. In z/OS V1R13, the limit is raised to 32 VLANs per stack per OSA port. There is no impact to OSA. This enhancement is in support of the emphasis on VLANs for network isolation in the zEnterprise IEDN.

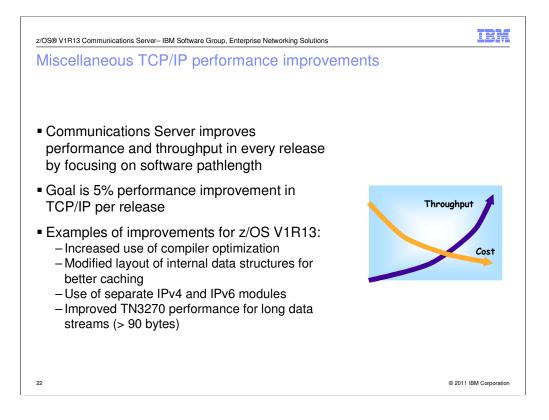
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Ado	ditional 64-b	bit exploitation		
• Tr	races relocat	ed to 64 bit common storage	Э	
-	-VTAM intern	al trace (VIT) moved from ECS	A and data space	
-	– Multiple CTF	ACE components moved from	data spaces to 6	4 bit common
ſ	CTRACE		z/IOS V1R13	
	Component	Current location	change	User
[	SYSTCPIP	TCPIPDS1 Dataspace	64 bit common	Stack
	SYSTCPIP	TN3270E Private storage	64 bit private	TN3270E
	SYSTCPDA	TCPIPDS1 Dataspace	64 bit common	Stack (NMI)
	SYSTCPIS	TCPIPDS1 Dataspace	64 bit common	Stack
	SYSTCPCN	TCPIPDS1 Dataspace	64 bit common	Stack (NMI only)
	SYSTCPSM	TCPIPDS1 Dataspace	64 bit common	Stack (NMI only)
	SYSTCPRE	Private SP229	No Change	RESOLVER
	SYSTCPRT	OMPROUTE Private storage	No Change	OMPROUTE
	SYSTCPIK	IKE daemon Private storage	No Change	IKESMP
	SYSTCPOT	TCPIPDS1 Dataspace	64 bit common	OSAENTA
	SYSTCPNS	NSS daemon's private storage	No Change	Security Server
0				© 2011 IBM Corpora

Multiple trace buffers have been relocated to take advantage of 64 bit common storage. The VTAM internal trace (VIT) is moved from ECSA to 64 bit common storage. In addition, the VIT data space is no longer used. This move is transparent to you if you use the external VIT to obtain trace records.

For TCP/IP, multiple CTRACE components have been moved from data spaces to 64 bit common storage. The table summarizes the changes. These moves are transparent to you as long as you use the NMI interface to obtain trace data.



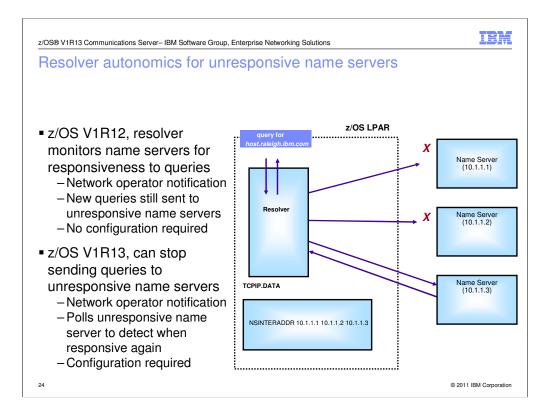
This slide shows a diagram of z/OS virtual storage. About 20 years ago, the focus was to move z/OS Communications Server storage utilization above the 31-bit "line" (16 MB). Now the attention is on moving storage utilization above the 64-bit "bar" (2 GB). In particular, reducing z/OS Communications Server's ECSA footprint has been an area of focus for several releases.



Communications Server development strives to improve performance and throughput in every release by focusing on software pathlength. The goal is at least a 5% performance improvement in TCP/IP per release. Some of the improvements for z/OS V1R13 are listed above.

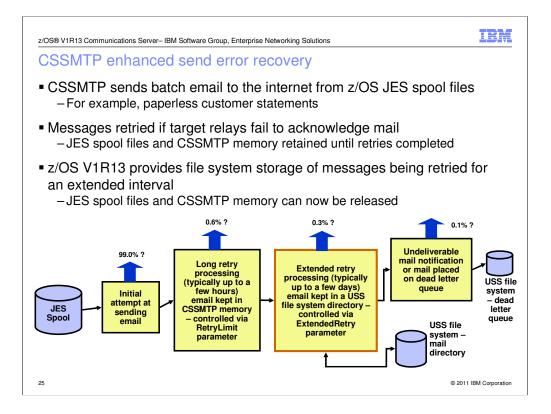
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The next topic is Availability enhancements.



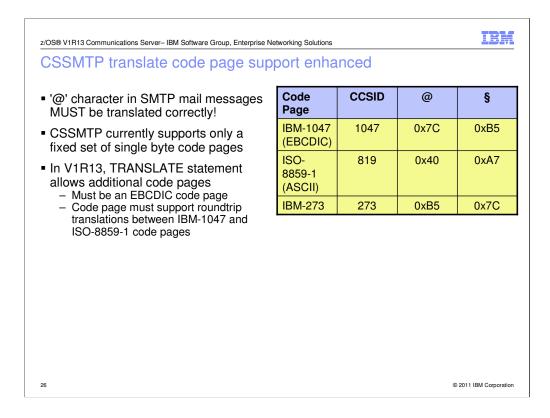
In z/OS V1R12, the resolver monitors name servers for responsiveness to queries. The network operator receives a notification when a name server becomes unresponsive. Although the resolver detected the unresponsive name server, new queries were still sent to that name server. This function is enabled by default.

In z/OS V1R13, the resolver can be configured to stop sending queries to unresponsive name servers. The resolver polls the unresponsive name server to detect when it becomes responsive again. The operator is notified of the condition using messages similar to those used in V1R12. The autonomic quiescing function must be explicitly enabled in the resolver setup file.

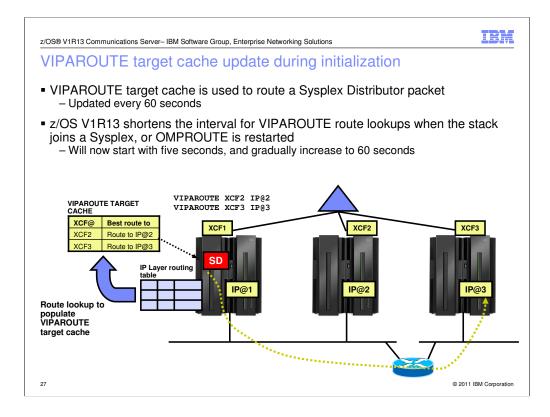


CSSMTP sends batch email to the internet from z/OS JES spool files. If the target relays fail to acknowledge mail, CSSMTP will retry for the configured interval up to five days then drop the message, and return an undeliverable notice. However, spool files cannot be deleted until all messages in the spool are delivered. A spool file can contain thousands of messages but only a few are being retried. Messages being retried are retained in CSSMTP memory.

z/OS V1R13 provides file system storage of messages being retried for an extended interval (beyond initial retry limit), so that JES spool files and CSSMTP memory can be released. It will continue to retry from memory and spool until the initial retry limit is reached. A new parameter can be configured to indicate how long beyond the existing interval to retry from the file system.



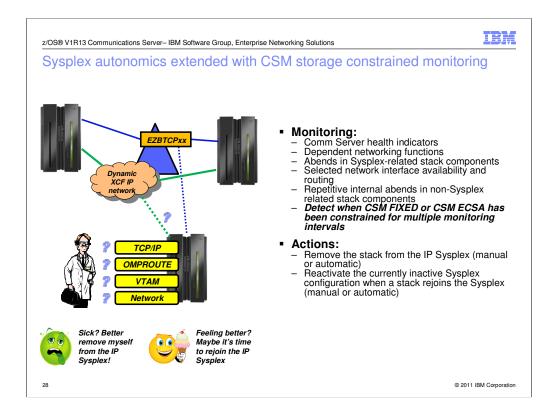
The '@' character has special meaning in SMTP mail messages. It must be translated correctly. CSSMTP currently supports only a fixed set of single byte code pages. Some installations use a code page that is not supported by the CSSMTP TRANSLATE statement. In z/OS V1R13, CSSMTP has been updated to allow the TRANSLATE statement to define additional code pages. Also, CSSMTP will allow the specification of a code page by its CCSID. Before z/OS V1R13, code pages must be a character string "IBM-XXXX", where XXXX is a subset of possible code pages. In z/OS V1R13, CSSMTP has been updated to expand the list of supported code pages. Also, CSSMTP will allow a user defined code page to be used. The code page must be an EBCDIC code page. The target ASCII code page is always ISO-8859-1 (or in reality, the US-ASCII subset of that code page). The code page must support roundtrip translations between the IBM-1047(EBCDIC) and the ISO-8859-1 (ASCII) code pages. The carriage return and line feed characters (CRLF), which are used to end the lines of commands and mail messages, must translate properly to ISO-8859-1 (x'0D0A').



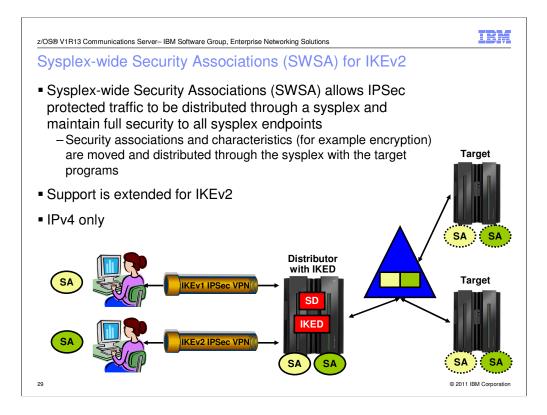
With VIPAROUTE support, distributed DVIPA traffic forwarded to a target stack can use any IP routing interface. An IP destination address is configured for each target XCF address. After the configuration profile is processed, normal IP routing rules are used to pick the best available route to that destination (IP routing table or policy-based routing). The IP route selected is cached for each target. Once a route to a target is cached, new connection requests distributed to that target use the cached target route. Periodically (every 60 seconds), a target cache refresh routine checks to see if the routing table has changed. If it has changed, a new route lookup is done to update each target's routing cache. If no route exists, the target is no longer considered as a distribution target; there is no fall back to use an XCF interface.

During a takeover that occurs because the primary routing stack is restarted, a route is unavailable until its interface finishes activating. If the initial target cache refresh occurs before this route is available, then there is a delay of close to 60 seconds until the next refresh. During this time, packets for distributed connections using VIPAROUTE might use non-optimal routes or might not be able to reach their targets.

This function substantially reduces this initial delay before acquiring a preferred VIPAROUTE by starting with smaller target cache refresh intervals. The initial refresh pattern is five, five, 15, 35, and 60 seconds. After reaching 60 seconds, the refresh interval will remain at 60 seconds. The initial refresh pattern is used whenever the stack joins the sysplex group or OMPROUTE is recycled.



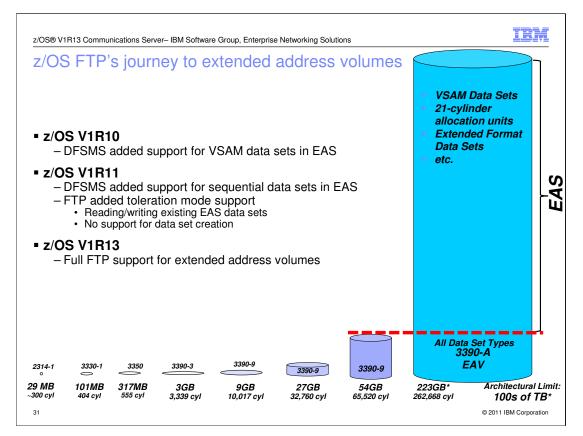
In V1R13, z/OS Communications Server more closely monitors common storage (CSM) usage and performs the recovery action if CSM is constrained (higher than 80% of defined CSM limit). Previous support only performs the recovery action when CSM is critical (higher than 90% of CSM limit). This reduces the impact of sysplex processing on other z/OS processes.



Sysplex Distributor negotiates security associations (SA) with the remote Client using the Internet Key Exchange protocol, IKE version 1 or IKE version 2. It sends copies of the SAs (shadows) to the targets for VPNs negotiated with either version of IKE. Targets use the SAs negotiated with either IKEv1 or IKEv2 to encrypt and decrypt data. The backup can recover SAs negotiated with either IKEv1 or IKEv1 or IKEv2 in case of DVIPA takeover. The coupling facility stores shared data for SAs negotiated with either IKEv1 or IKEv2.

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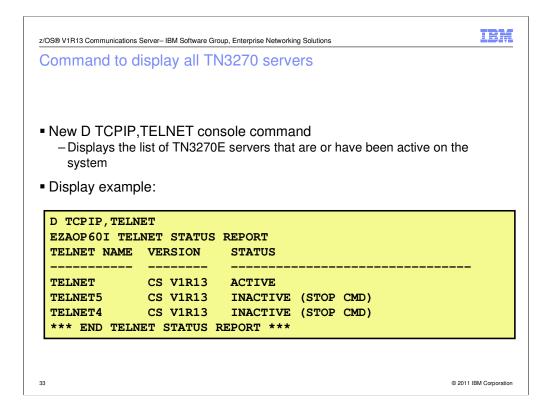
The next topic is Application, Middleware and Workload Enablement enhancements.



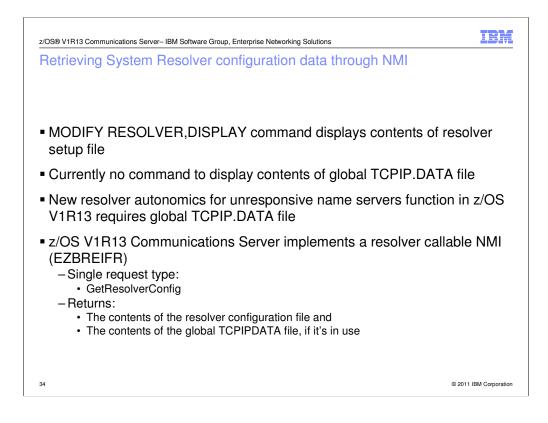
In z/OS V1R10, DFSMS added support for VSAM data sets in the extended addressing space (EAS). FTP does not support VSAM data sets, so there was no impact. In z/OS V1R11, DFSMS added support for extended format sequential data sets eligible to reside in the EAS. FTP added support for reading and writing existing EAS data sets, but not for creating them (toleration mode). In z/OS V1R13, FTP provides full support for extended address volumes (EAV) without requiring use of an SMS data class. Before z/OS V1R13, you can use an EAV for SMS managed datasets if set up properly in an SMS data class.

	BASIC FORMAT	LARGE FORMAT	EXTENDED FORMAT
DSNTYPE parameter value	BASIC	LARGE	EXTPREF
Maximum tracks per volume	65,535	16,777,215	EXTREQ Larger than 16,777,215
Maximum extents per volume	16	16	123
Why choose this format?	Maximum compatibility.	Can be much larger than basic format.	Can be much larger than basic format. Can be striped, compressed format or an combination.
z/OS V1R13	support for large fo large format data so port for large forma	ets completed	

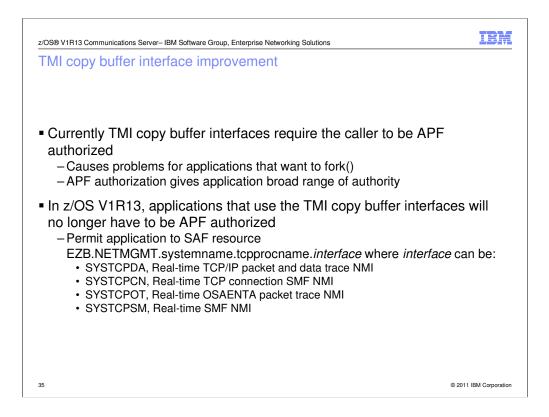
Large format physical sequential datasets can have more than 64K tracks per volume. They do not have to. They can have fewer than 64K tracks per volume and still be large *format*. DFSMS has offered large format physical sequential datasets since z/OS V1R7. The access methods supported are BSAM, QSAM, and EXCP. Language Environment (LE) Runtime Library large format dataset support completed in z/OS V1R13 now enables z/OS FTP support for these datasets. FTP support means sending from large format data sets and allocating and storing into large format data sets. In block mode, you can use the restart subcommand to resume the failed transfer to and from large format data sets. Support is also added for transfers to and from z/OS UNIX files larger than two gigabytes.



A new D TCPIP, TELNET console command has been added to display the list of TN3270E servers that are or have been active on the system. This can be a starting point for performing automation on TN3270E servers. After all, you have to know what's there before you can operate on it!

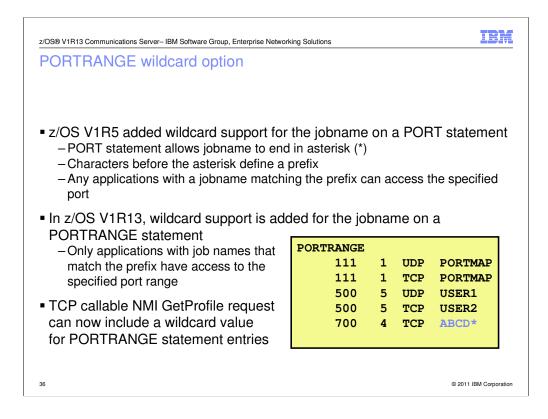


You can use the MODIFY RESOLVER, DISPLAY command to display the contents of the resolver setup file. However, there is no resolver command available to display the contents of the global TCPIP.DATA file. Before z/OS V1R13, you must collect Trace Resolver output to see the global TCPIP.DATA settings dynamically. The new resolver function in z/OS V1R13 to monitor and quiesce unresponsive name servers, depends on a global TCPIP.DATA file. Only name servers specified in this global TCPIP.DATA file are monitored. z/OS V1R13 Communications Server implements a resolver callable NMI (EZBREIFR). The NMI provides a high-speed, low-overhead callable programming interface for network management applications to access data related to the resolver. There is a single request type (GetResolverConfig) which returns the contents of the resolver configuration file and the contents of the global TCPIP.DATA file, if it is in use.



The TCP/IP Management Interface (TMI) copy buffer interfaces currently require the caller to be APF authorized. The interfaces are EZBTMIC1, EZBTMIC4, and TMI\_Copybuffer(). This causes problems for applications that want to fork() as APF authorization is not inherited. In addition, APF authorization gives applications a broad range of authority.

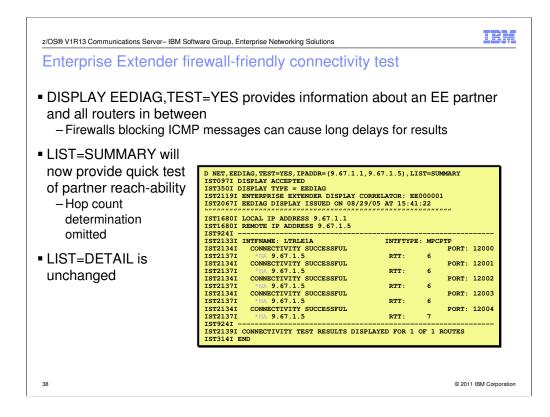
In z/OS V1R13, network management applications that use these TMI copy buffer interfaces will no longer have to be APF authorized. As an alternative to APF authorization, the user ID that the application is running under can be authorized to the appropriate SAF resource EZB.NETMGMT.systemname.tcpprocname.interface. The interface can be SYSTCPDA, SYSTCPCN, SYSTCPOT, or SYSTCPSM.



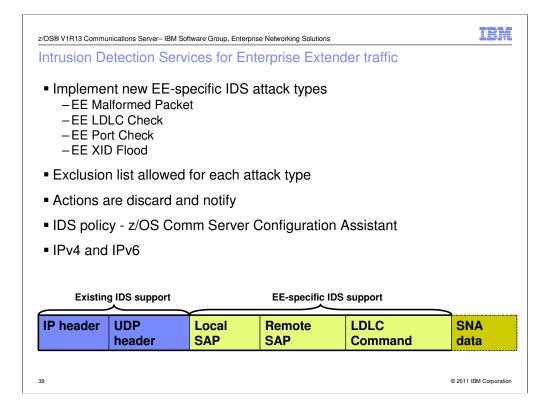
A wildcard option has been allowed for the job name value specified on a PORT statement since z/OS V1R5. The PORT statement allows a job name to end in an asterisk. The characters before the asterisk define a prefix. Any applications with a job name matching the prefix can access the specified port. Similar wildcard support was not provided for the PORTRANGE statement. The PORTRANGE statement in z/OS V1R13 allows a wildcard job name. As with the PORT statement, the job name can end in an asterisk. The characters before the asterisk define a prefix. Only applications with job names that match the prefix have access to the specified port range. As a result of this, the GetProfile request of the TCP callable NMI (EZBNMIFR) can now include a wildcard value in the NMTP\_PORTJobName field for entries that represent a PORTRANGE statement.

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The next topic is Enterprise Extender and SNA enhancements.



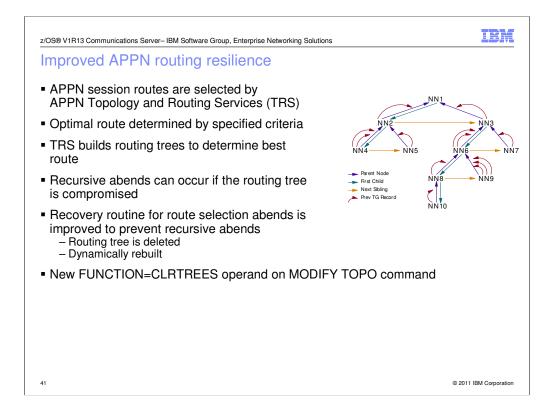
The DISPLAY EE DIAG TEST=YES command provides information about an Enterprise Extender partner and all of the routers in between. If a firewall in between is blocking ICMP messages, there can be a long delay before getting results. The delay is because of timeouts waiting for ICMP messages that never come. The delay is the number of router hops past the firewall times nine seconds. The existing LIST=SUMMARY option now provides a quick test of partner reachability. The probe is sent to the partner with TTL=255 so it does not probe any intermediate hops. The output no longer includes the hop count determination. The LIST=DETAIL output is unchanged and includes intermediate hop information and the hop count determination.



Intrusion detection services (IDS) is enhanced to implement four new IDS attack types for Enterprise Extender (EE). These attack types are supported for IPv4 and IPv6 EE traffic. The EE Malformed Packet attack type checks inbound EE packets for incorrect lengths. The EE LDLC Check attack type checks that inbound LDLC control commands are only received on the signaling port (12000). The EE Port Check attack type checks that inbound EE packets contain matching source and destination ports. The EE XID Flood attack type checks if a threshold is met for inbound XIDs within one minute. The actions allowed are to discard the packet and to provide a notification. The EE XID Flood attack only supports the notify action. An exclusion list can be created to exclude specific hosts from attack checking. Events notifications can be sent to syslogd, to the console, to IDS packet trace and to Tivoli Security Operations Manager (TSOM). IDS is configured using policies and is supported by the Configuration Assistant.

		uration Assistan erprise Extender			
New Requirement	nt attack protection —				
3. To enable protection for a spe You will be prompted for additi Attack type	cific attack type, selec	w from the Enabled protection table a a row from the Attack type table and your attack type selection. Fill in the of Enabled protection	click the "Enable" button.	1.00	
Data Hiding Attack Flood Attack Global TCP Stall Attack ICMP Redirect Attack IPv4 Options Attack IPv4 Options Attack IPv4 Optionund Raw Attack	Enable>	EE LDLC Check Attack	EEMalformedPacket EELDLCCheck EEPortCheck EEXIDFlood	Action Report Events Drop Packets or Connection Both Drop and Report Report Events	> > > > > >
IPv4 Protocols Attack IPv6 Destination Options Attack IPv6 Hop-by-Hop Options Attack IPv6 Next Header Attack IPv6 Outbound Raw Attack	<ul> <li>✓</li> </ul>				
Default Report Settings for Attack	s	Modify Copy	Advanced View Deta	ils)	
Help ?			< <u>B</u> ack	Next > Einish	Cance

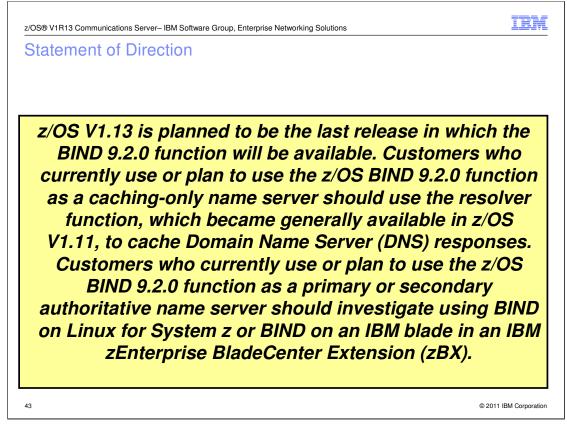
The Configuration Assistant was updated to support the four new EE attack types. The window here shows a new requirement map with the four EE attack types listed. You can disable an attack type by selecting it and clicking the disable button.



APPN session routes are selected by APPN Topology and Routing Services (TRS). These routes are used to locate resources (directed search routes) and for LU-LU sessions. The optimal route is determined by specified criteria such as line speed, cost of data transmitted, security, and user-defined values. TRS builds complex routing trees to determine the best path. The tree records represent nodes along preferred routes. Recursive abends can occur if a pointer in a routing tree is compromised. Every time a session route is requested using that tree, abends occur and the session fails. VTAM must be restarted to recover. The recovery routine for route selection abends is improved in z/OS V1R13 to prevent recursive abends – by removing the entire storage area that contains the routing tree and dynamically rebuilding it. Existing sessions are not affected. It can temporarily impact VTAM performance in large APPN networks as the routing tree is being rebuilt. For the rare case where the routing tree is corrupted, but no abends occur, a new command can be used to perform the same process as described above.

z/OS® V1R13 Communications Server- IBM Software Group, Enterprise Networking Solutions	IEM
Agenda	
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Security	
Simplification	
Economics and Platform Efficiency	
Availability	
Application/Middleware/Workload Enablement	
EE/SNA	
SOD and Reference	
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The next topic is Enterprise Extender and SNA enhancements.



The next topic is Enterprise Extender and SNA enhancements.

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http://www.ibm.com/systems/z/os/zos/bkserv/	IBM z/OS Internet library – PDF files of all z/OS manuals including Communications Server

This page provides a series of URLs for web pages where additional information on z/OS Communications Server can be found.